# Dynamic Programming IV: Shortest Paths

## The Problem

Given a graph $G = (V, E)$ with *any edge weights* $w_e$ for all $e \in E$, a source $s$, output the shortest $s - v$ path for every $v \in V$.

## Naïve Approach

Let $N(v)$ be the neighbors of $v$.

- Subproblem: Let $\text{OPT}(v)$ denote the length of the shortest path from $s$ to $v$.

- Recurrence: $\text{OPT}(v) = \min_{u \in N(v)} \text{OPT}(u) + w_{(u,v)}$.

- Base Case: $\text{OPT}(s) = 0$.

What's wrong with this approach? List any issues you see:

# The Bellman-Ford Algorithm

Main Idea: Impose a measure of progress—parametrize the subproblems.

More specifically:

- Consider a shortest path from $s \to u \to v \to t$ with $k$ edges

- $s \to u \to v$ is a shortest $s \to v$ path with $k - 1$ edges

Measure of progress: the number of edges in path. First compute all shortest paths with $\leq 1$ edges. Then shortest paths with $\leq 2$ edges. And so on, until we compute shortest paths of length $\leq n - 1$.

**Structural Observation:**

*Proof.*

**Formal Description**

**State Your Subproblem:**

**State Your Recurrence:**

**Prove Your Recurrence:**

**State Your Base Cases:**

**Present Your Algorithm:**

**Runtime:**

- Size of table:

- Time to fill in table:

    (How many table lookups?)

Total time:

**Better Runtime Analysis**

- How many row updates?

- How many table lookups per row update?

- Runtime =

**Detecting Negative Cycles**

Idea: Add an extra column (Column $n$) to the memo table.

**Claim 1.** There is a negative cycle in the graph if and only if Column $n \neq$ Column $n - 1$.

*Proof.* ($\Longleftarrow$) Suppose the graph has no negative cycles. Then by the Structural Observation,

($\Longrightarrow$) Suppose there is a negative cycle in the graph.

To detect negative cycles:

- Run Bellman-Ford, add an $n$th column.
- If Column $n =$ Column $n - 1$, return "No Negative Cycles"
- Otherwise, return that there's a negative cycle.

## Space-Efficient Bellman-Ford

How much space does Bellman-Ford use?

**Observation:**

**Improved algorithm:**

Total space usage: